

Comparison of Symbolic Maximal End Component Decomposition Algorithms

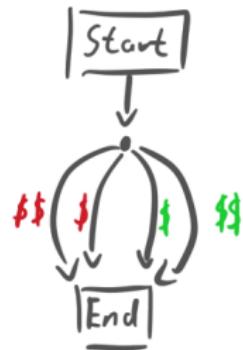
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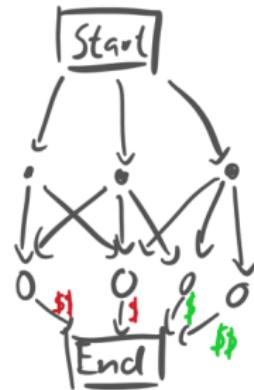
12.09.2023

(<https://FelixFaber.dev/> for Thesis / Slides)

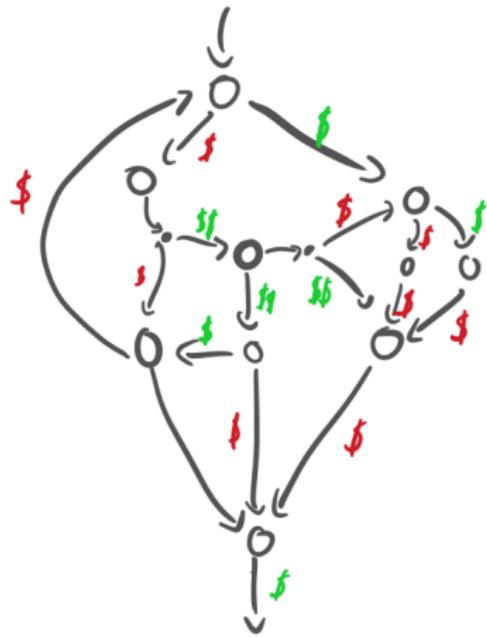
Motivation



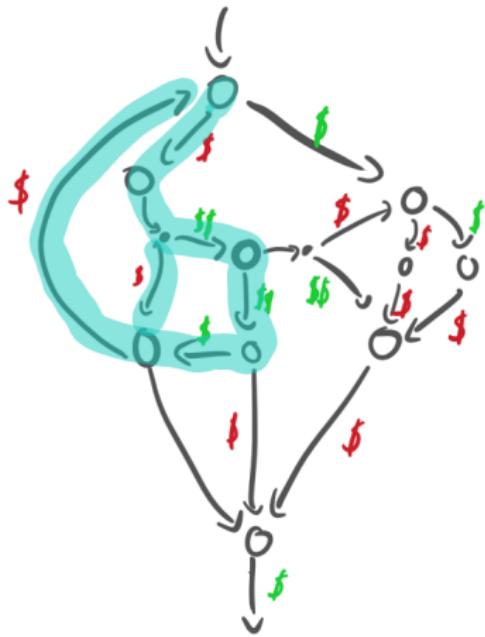
Motivation



Motivation

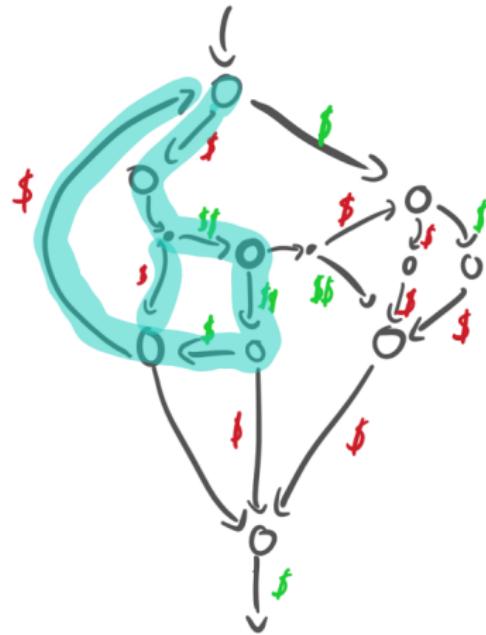


Motivation



Motivation

- ▶ Represent slot machine as model
- ▶ Find all “loops”
- ▶ Check for exploits
- ⇒ (Symbolic) model checking



Symbolic MEC Decomposition

$$n := |V|, \quad m := |E|$$

Algorithm	Worst-Case Sym. Ops.	Worst-Case Sym. Space	Applicability
NAIVE [DA98, CHL ⁺ 18]	$O(n^2)$	$O(\log n)$	G_{EBA}, G_{VBA}
LOCKSTEP [CHL ⁺ 18]	$O(n\sqrt{m})$	$O(\sqrt{m})$	G_{EBA}, G_{VBA}
COLLAPSING [CDHS21]	$O(n^{2-\epsilon} \log n)$ $(0 < \epsilon \leq 0.5)$	$O(n^\epsilon \log n)$ $(0 < \epsilon \leq 0.5)$	G_{VBA}

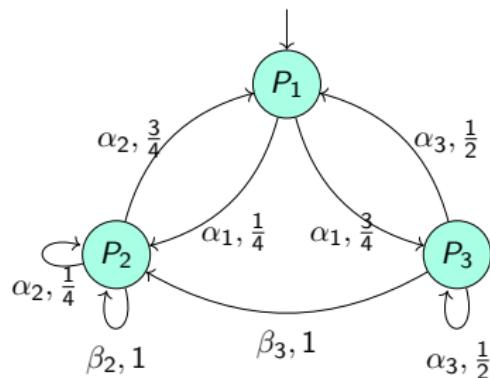
⇒ Performance in practice?

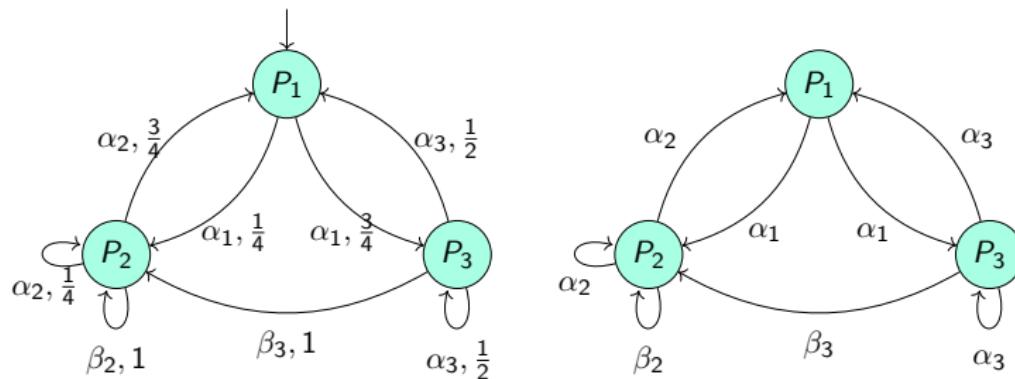
Structure

- ▶ MDPs: MECs on graph-like structures G_{EBA} and G_{VBA}
- ▶ Symbolic MDP representation using BDDs
- ▶ Symbolic MEC algorithms: NAIVE, LOCKSTEP, COLLAPSING
- ▶ Empirical evaluation of algorithms
- ▶ Summary and Future Work

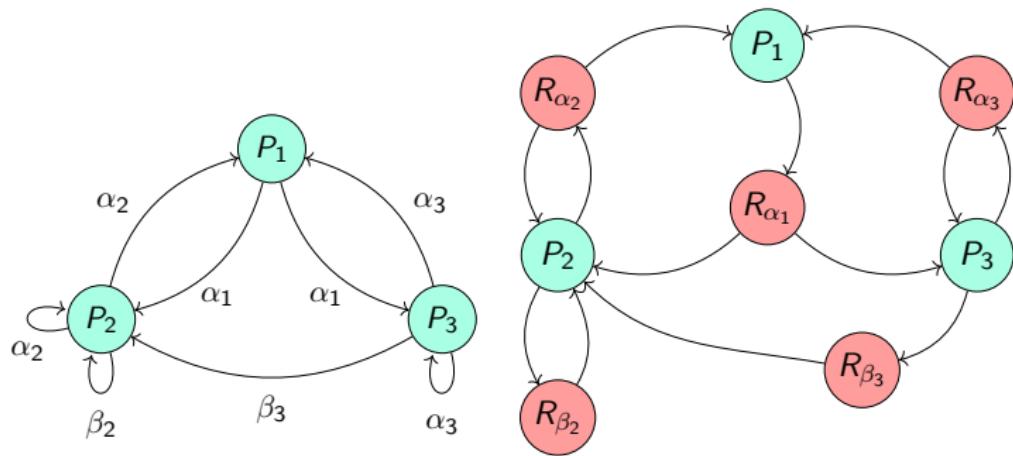
MDP

$$\mathcal{M} = (S, A, d_{\text{init}}, \delta, r)$$



Edge-based actions G_{EBA} 

Edge-based actions G_{EBA}
vs
Vertex-based actions G_{VBA}



ECs

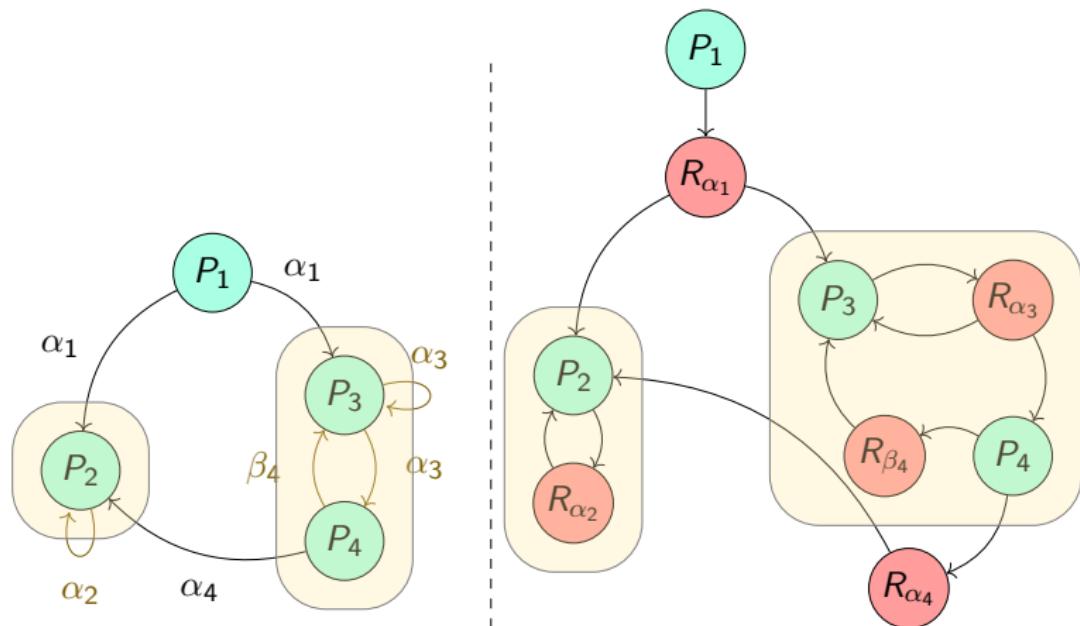
End Component:

- ▶ Set of states and actions
- ▶ Each state has at least one action
- ▶ All pairs of states of EC are reachable from another (using actions of EC)
- ▶ No outgoing actions
 - ⇒ Player can stay within EC indefinitely

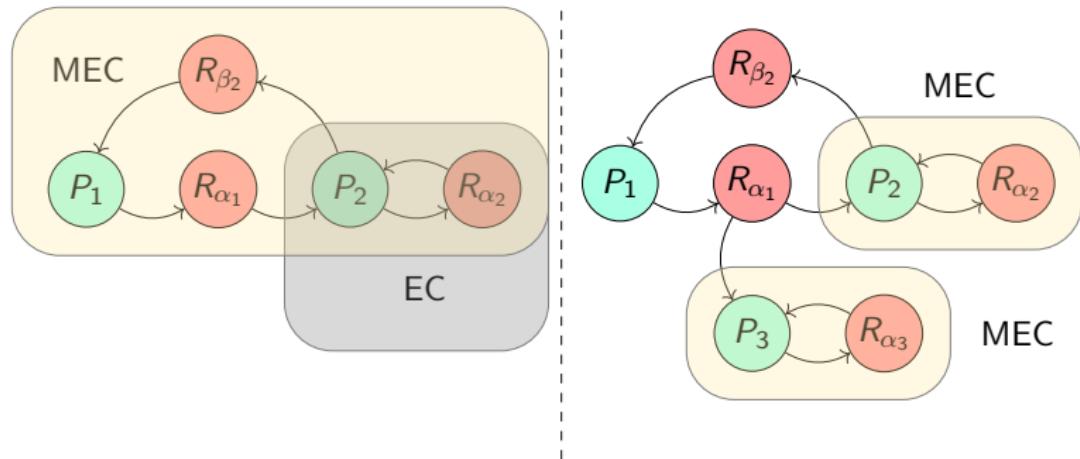
MECs

- ▶ MEC: Maximal EC (WRT set inclusion)
 - ⇒ Each state/action belongs to at most one MEC
- ▶ MEC decomposition: Compute all MECs of MDP \mathcal{M}

MECs



MECs

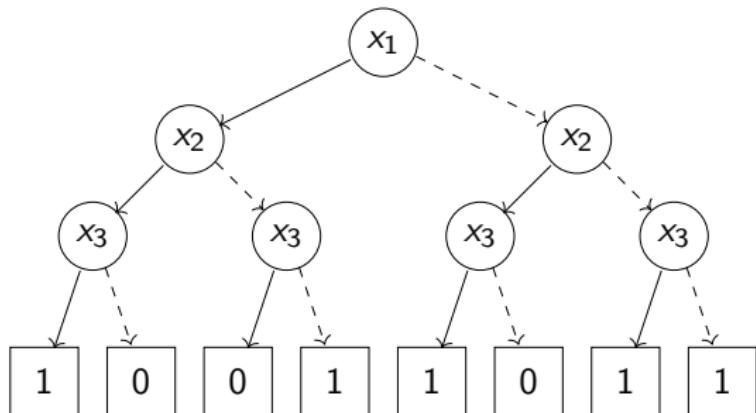


BDDs

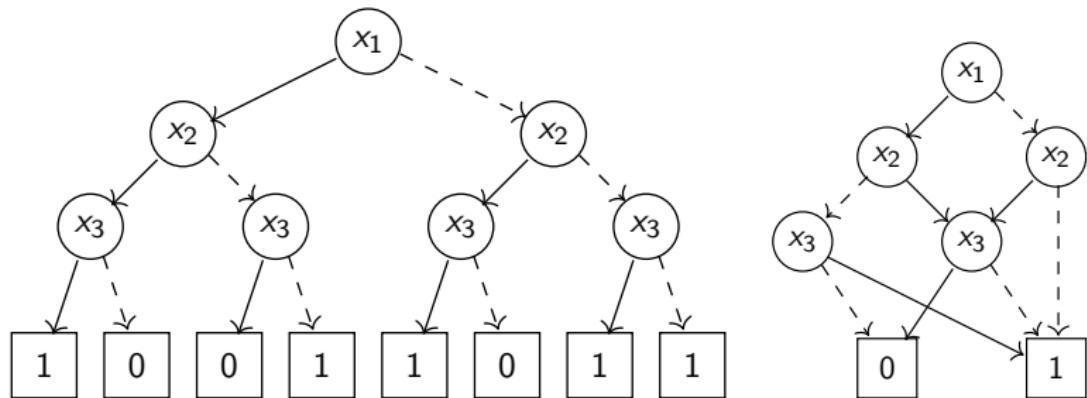
Binary Decision Diagrams (BDDs)

In practice: Reduced Ordered BDDs (ROBDDs)

x_1	x_2	x_3	f
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1



BDDs



MDP Representation: States

$$f_{V'} : \{0, 1\}^t \rightarrow \{0, 1\}, \quad f_{V'}(\underbrace{x_1, \dots, x_t}_{\text{encodes } s \in V}) = \begin{cases} 1, & s \in V' \\ 0, & s \notin V' \end{cases}$$

G_{VBA} :

- ▶ f_{V_P} for Player Vertices V_P
- ▶ f_{V_R} for Random Vertices V_R

MDP Representation: Transitions

Transition BDD G_{VBA} :

$$t_{VBA}(\underbrace{x_1, \dots, x_t}_{\substack{s \in V \\ (\text{Row})}}, \underbrace{x'_1, \dots, x'_t}_{\substack{s' \in V \\ (\text{Column})}}) = \begin{cases} 1, & (s, s') \in E \\ 0, & (s, s') \notin E \end{cases}$$

Transition BDD G_{EBA} :

$$t_{EBA}(\underbrace{x_1, \dots, x_t}_{\substack{s \in V \\ (\text{Row-Group})}}, \underbrace{y_1, \dots, y_u}_{\alpha \in A[s]}, \underbrace{x'_1, \dots, x'_t}_{\substack{s' \in V \\ (\text{Column})}}) = \begin{cases} 1, & (s, \alpha, s') \in E \\ 0, & (s, \alpha, s') \notin E \end{cases}$$

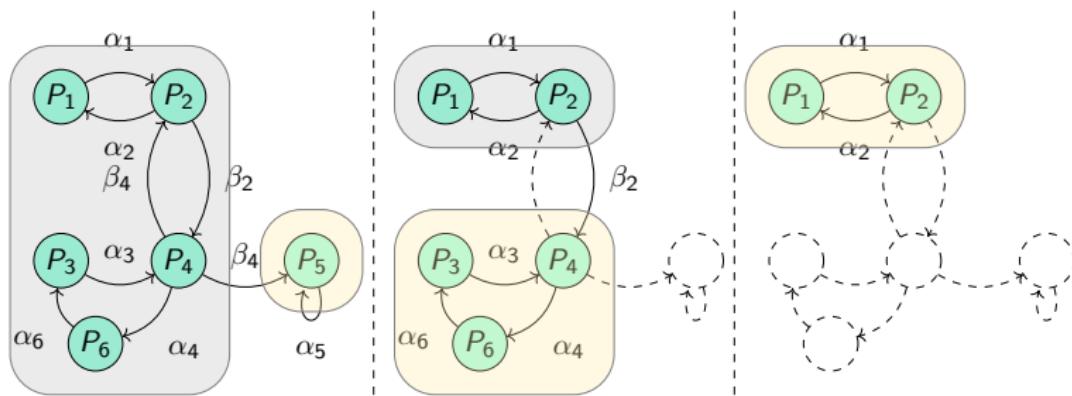
Symbolic Algorithms

- ▶ Need to utilize symbolic operations
- ▶ Runtime: symbolic operations (*Pre/Post*)
- ▶ Space: symbolic space (BDD)
- ▶ Assumes $O(n)$ symbolic SCC decomposition [GPP03]

$$n := |V|, \quad m := |E|$$

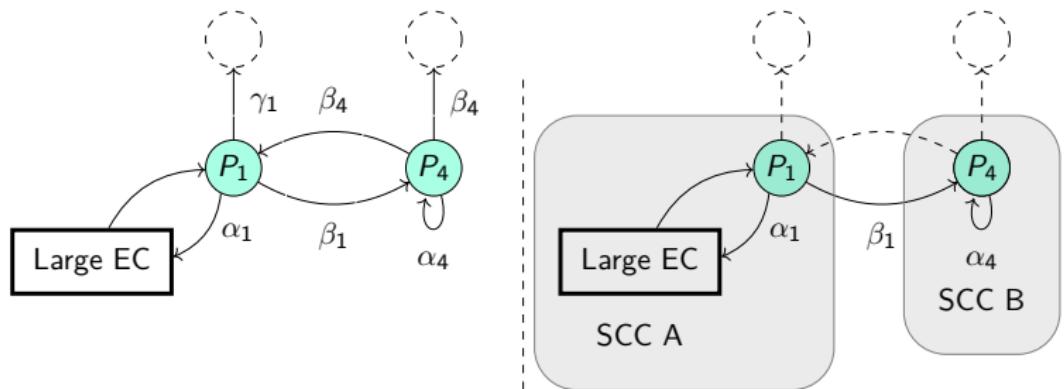
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MEC Decomposition: NAIVE



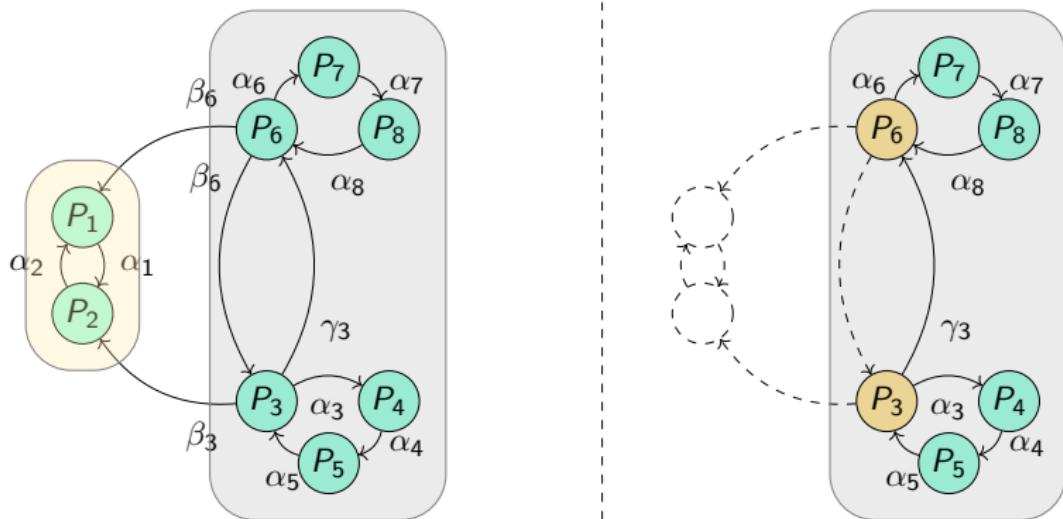
Symbolic implementation of basic MEC decomposition algorithm
⇒ SCC decompositions and removal of outgoing actions

MEC Decomposition: LOCKSTEP



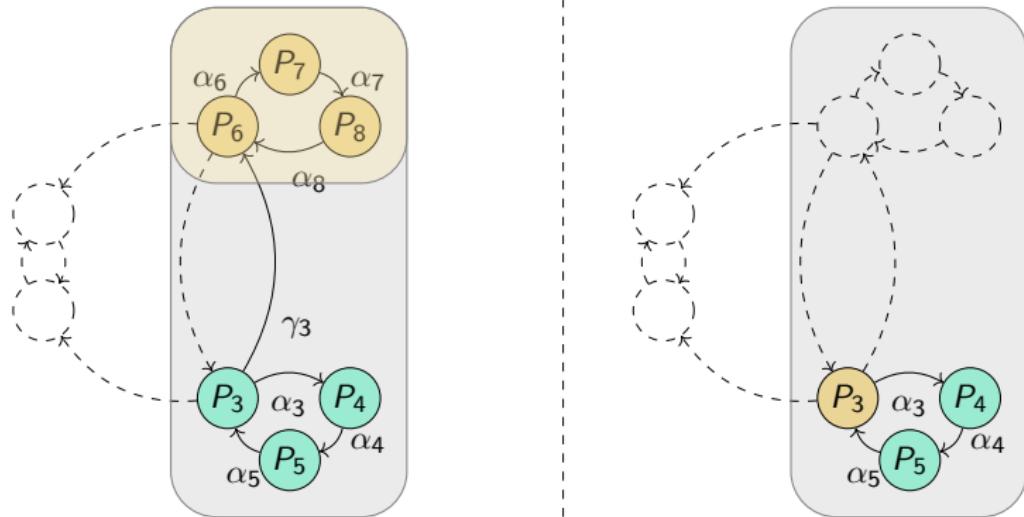
Idea: identify bottom SCCs to avoid duplicate vertex traversals

MEC Decomposition: LOCKSTEP

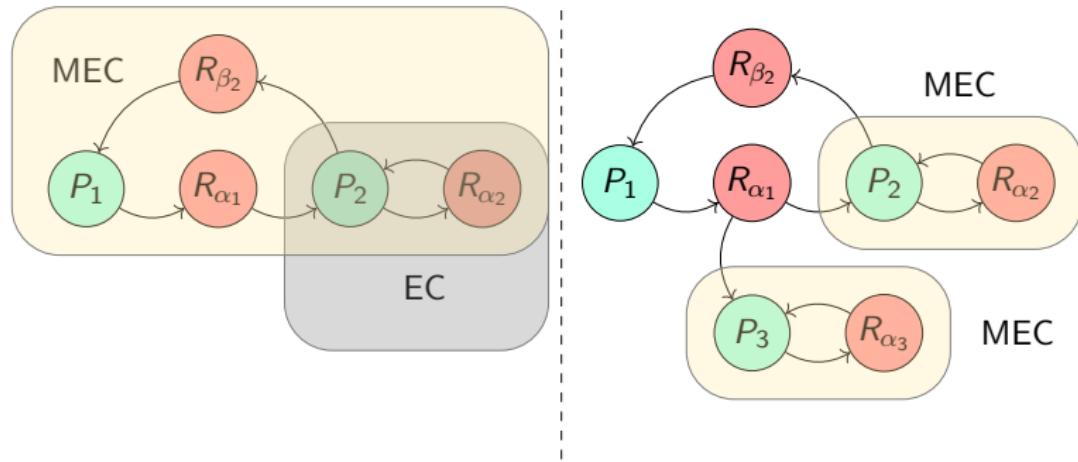


- ▶ Start lockstep search from each vertex which lost an edge.
- ▶ Post operations yield bottom SCC

MEC Decomposition: LOCKSTEP



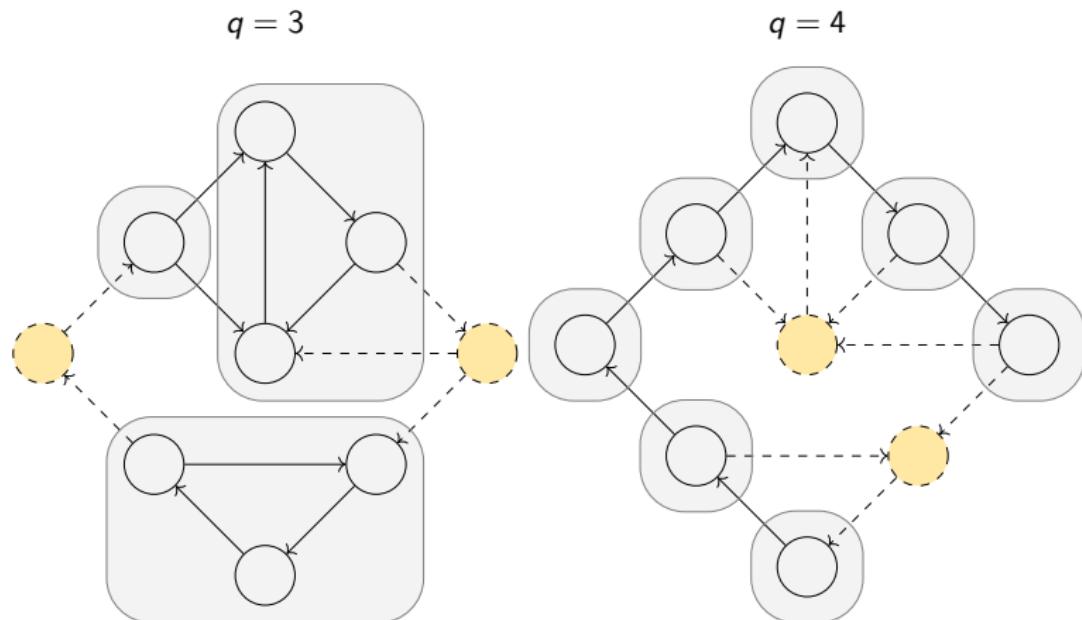
MEC Decomposition: COLLAPSING



Two passes:

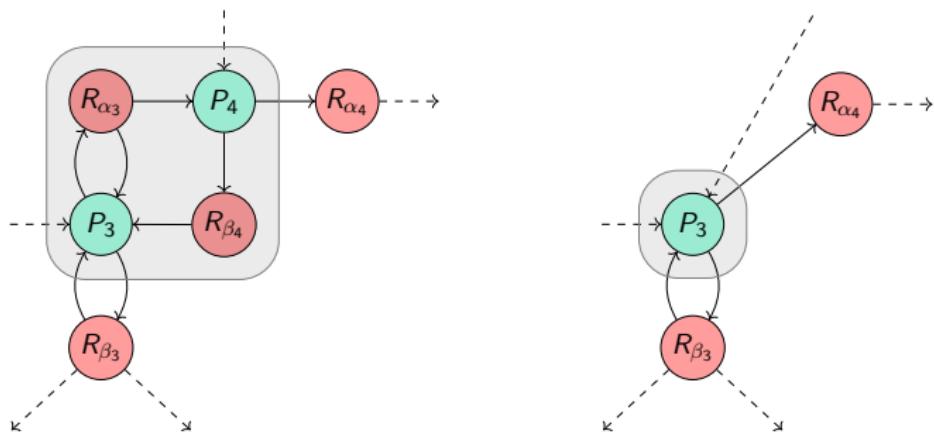
- ▶ Detect vertices of ECs quickly
- ▶ SCC decomposition on ECs \Rightarrow MEC decomposition

MEC Decomposition: COLLAPSING



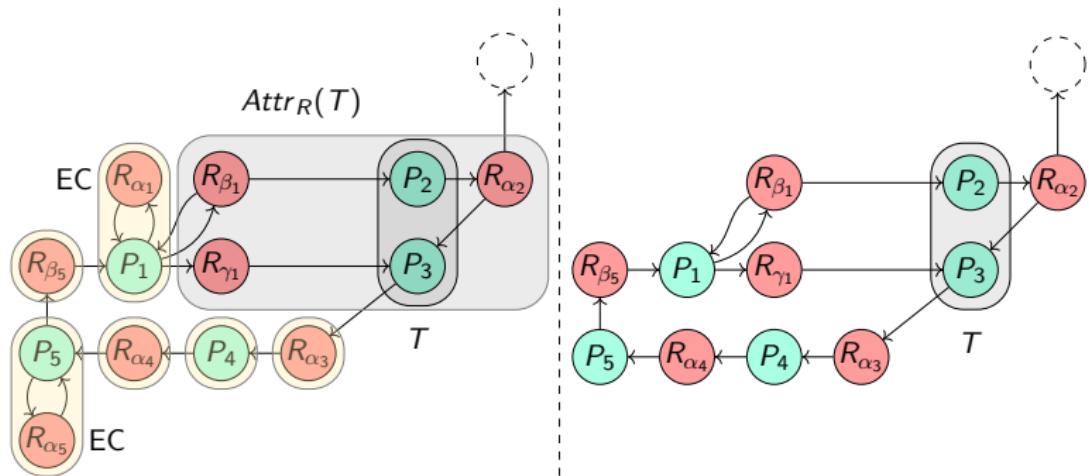
- ▶ Split up large SCCs using separator.
- ▶ Process smaller SCCs
- ▶ Re-add separator

MEC Decomposition: COLLAPSING

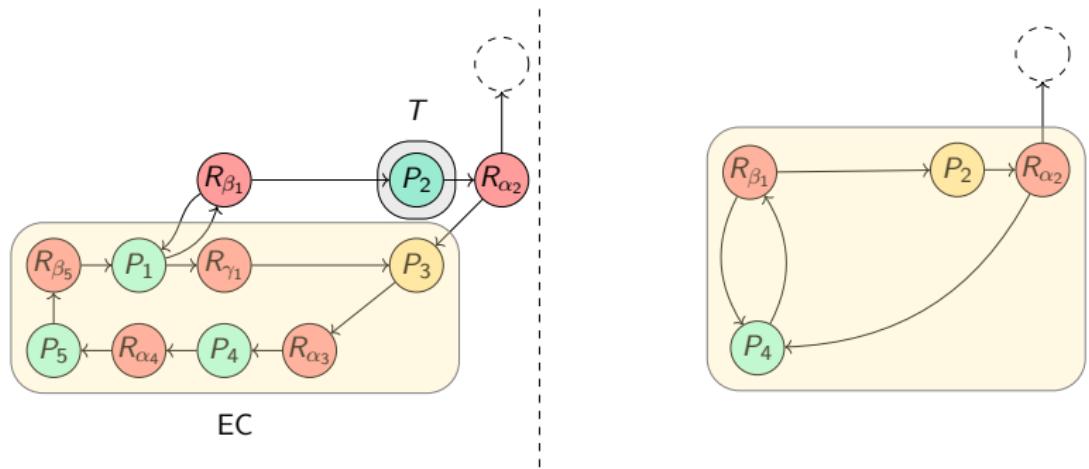


Retraversal of ECs cheaper after collapsing

MEC Decomposition: COLLAPSING

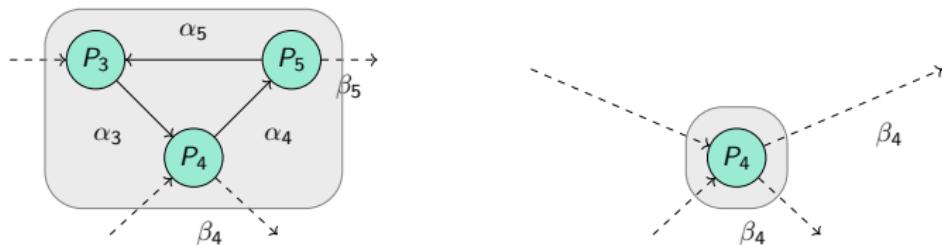


MEC Decomposition: COLLAPSING



MEC Decomposition: COLLAPSING

$$t_{EBA}(\underbrace{x_1, \dots, x_t}_{\substack{\text{Source Vertex} \\ s \in V}}, \underbrace{y_1, \dots, y_u}_{\substack{\text{Action} \\ \alpha \in A[s]}}, \underbrace{x'_1, \dots, x'_t}_{\substack{\text{Destination Vertex} \\ s' \in V}}) = \begin{cases} 1, & (s, \alpha, s') \in E \\ 0, & (s, \alpha, s') \notin E \end{cases}$$



Evaluation

$$n := |V|, \quad m := |E|$$

Algorithm	Worst-Case Sym. Ops.	Worst-Case Sym. Space	Applicability
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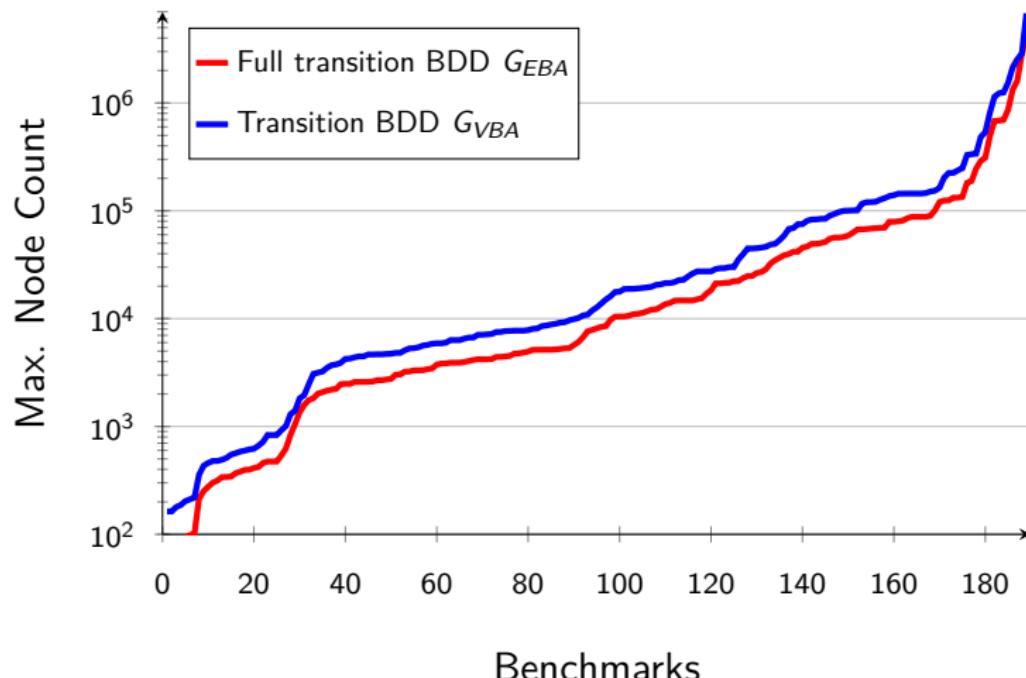
Performance in practice?

→ Empirical evaluation using

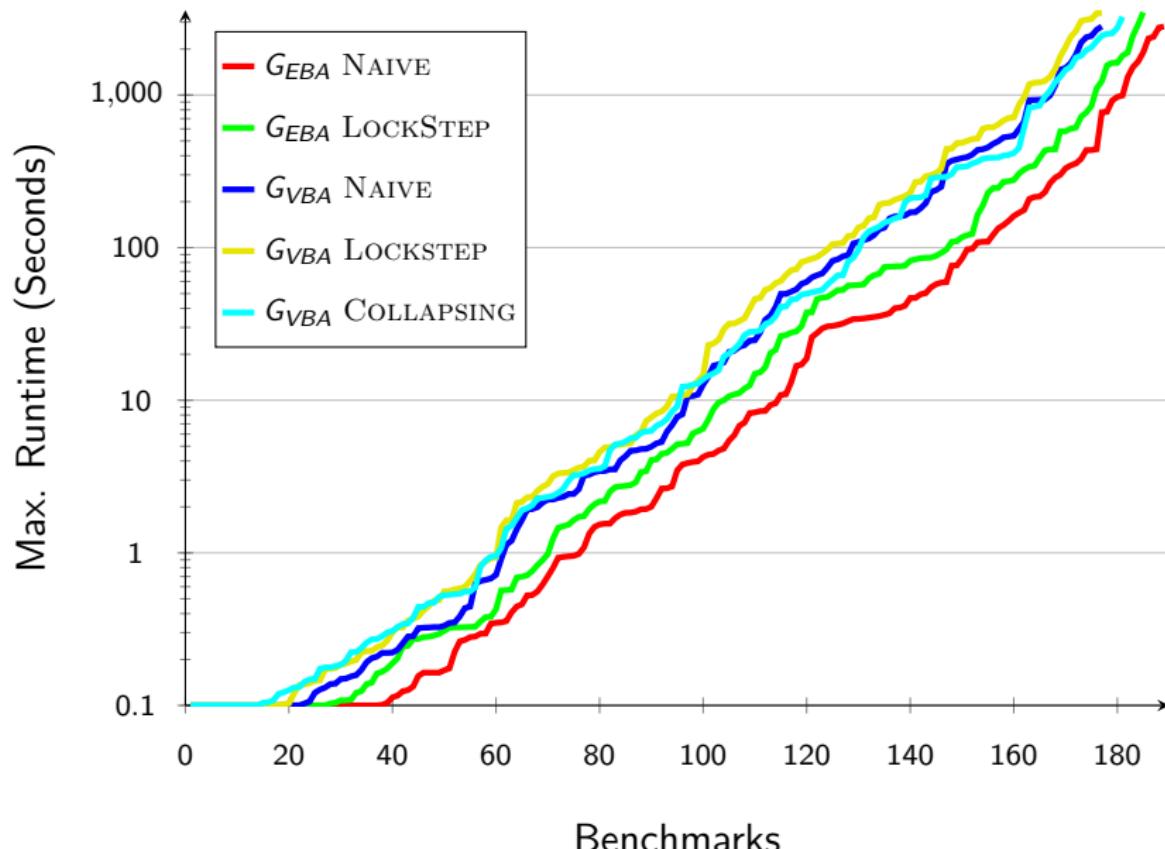
- ▶ STORM
- ▶ quantitative verification benchmark set (QVBS)
- ▶ 1 Hour time limit

Evaluation

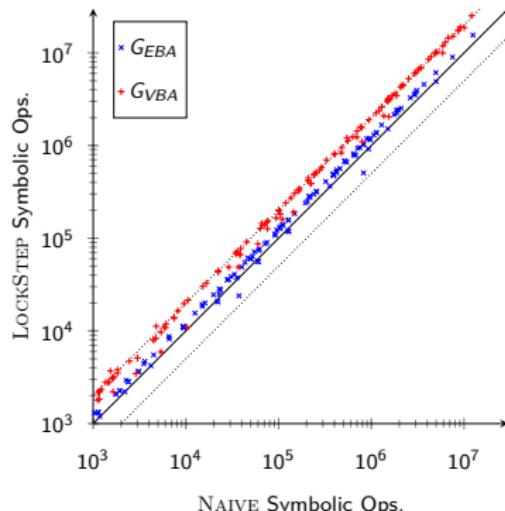
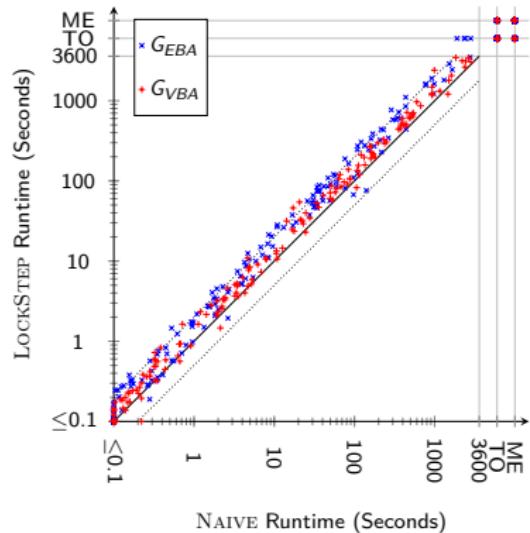
- ▶ STORM assumes G_{EBA} \Rightarrow convert to G_{VBA} if desired
- ▶ 189 Benchmarks for runtimes
- ▶ 177 Benchmarks for symbolic operations



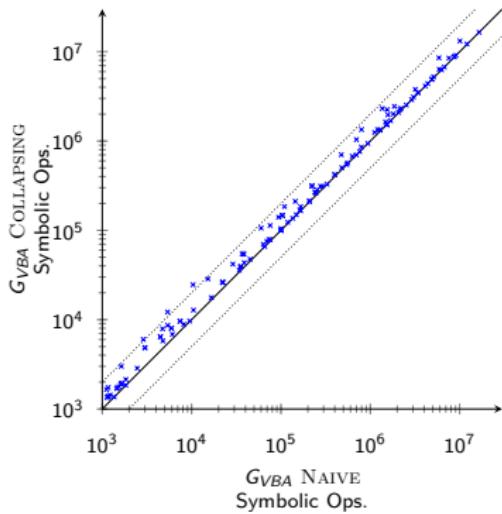
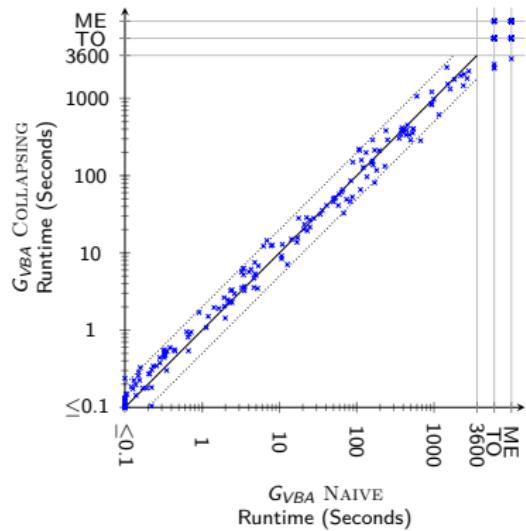
Evaluation: Runtime Overview



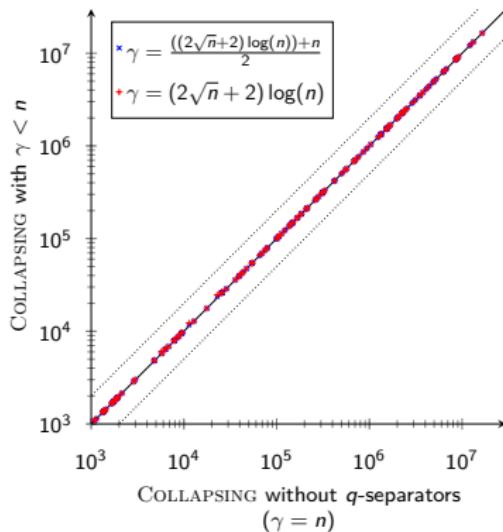
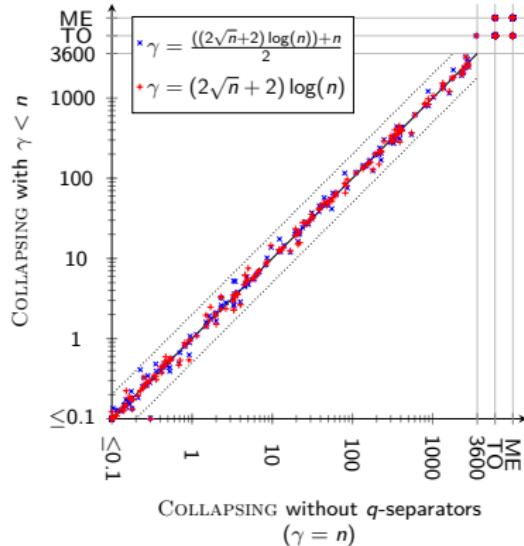
Evaluation: LOCKSTEP



Evaluation: COLLAPSING



Evaluation: COLLAPSING



Summary

- ▶ Theoretical worst-case improvements are not reflected in empirical results
- ▶ Least amount of Sym. Ops.: NAIVE
- ▶ Best Runtime: NAIVE or COLLAPSING
 - ▶ But COLLAPSING is (at the moment) only applicable to G_{VBA}

Future Work

- ▶ “Native” G_{VBA} implementation
- ▶ COLLAPSING implementation on G_{EBA}
- ▶ Improvements in SCC decomposition [LSS⁺23]
- ▶ Comparison of *explicit* MEC decomposition algorithms

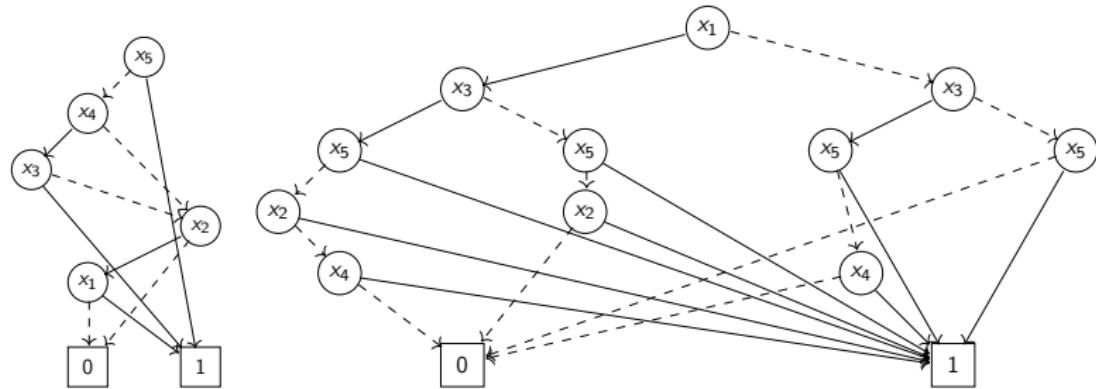
References I

- [CDHS21] Krishnendu Chatterjee, Wolfgang Dvořák, Monika Henzinger, and Alexander Svozil. Symbolic time and space tradeoffs for probabilistic verification. In *2021 36th Annual ACM/IEEE Symposium on Logic in Computer Science (LICS)*, pages 1–13. IEEE, 2021.
- [CHL⁺18] Krishnendu Chatterjee, Monika Henzinger, Veronika Loitzenbauer, Simin Oraee, and Viktor Toman. Symbolic algorithms for graphs and Markov decision processes with fairness objectives. In *International Conference on Computer Aided Verification*, pages 178–197. Springer, 2018.
- [DA98] Luca De Alfaro. *Formal verification of probabilistic systems*. stanford university, 1998.

References II

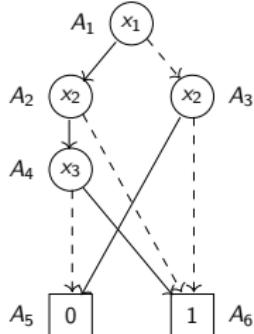
- [GPP03] Raffaella Gentilini, Carla Piazza, and Alberto Policriti. Computing strongly connected components in a linear number of symbolic steps. In *SODA*, volume 3, pages 573–582, 2003.
- [LSS⁺23] Casper Abild Larsen, Simon Meldahl Schmidt, Jesper Steensgaard, Anna Blume Jakobsen, Jaco van de Pol, and Andreas Pavlogiannis. A Truly Symbolic Linear-Time Algorithm for SCC Decomposition. In *International Conference on Tools and Algorithms for the Construction and Analysis of Systems*, pages 353–371. Springer, 2023.

ROBDDs: Variable Ordering

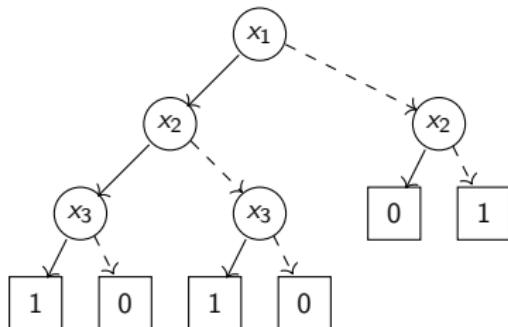
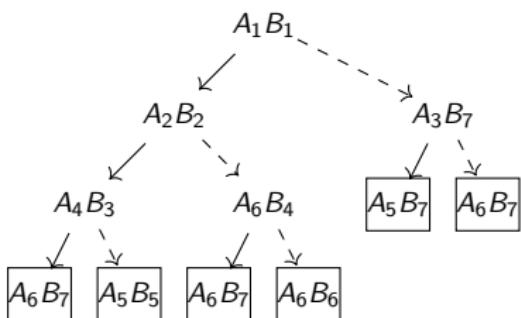
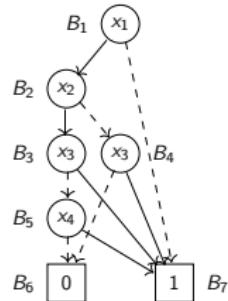


ROBDDs: AND Operation

$$f(x_1, \dots, x_4) = x_1x_3 + \bar{x}_2$$

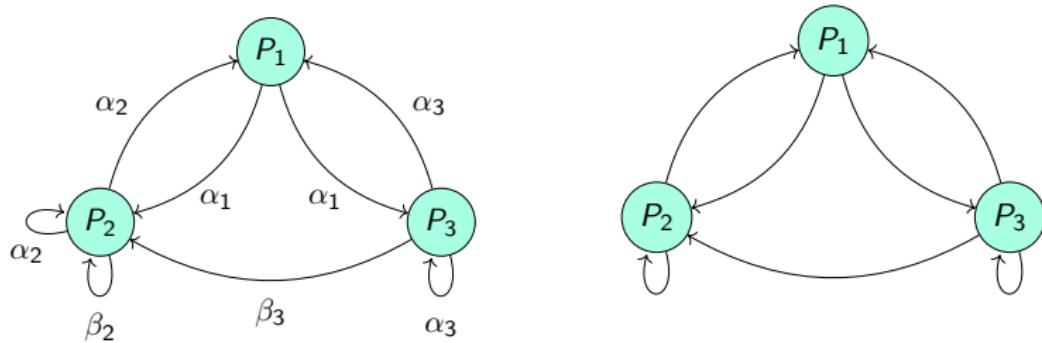


$$g(x_1, \dots, x_4) = x_2x_4 + x_3 + \bar{x}_1$$



MDP: Directed Graph of G_{EBA}

Directed graph of G_{EBA}



MDP Representation: Explicit Transitions

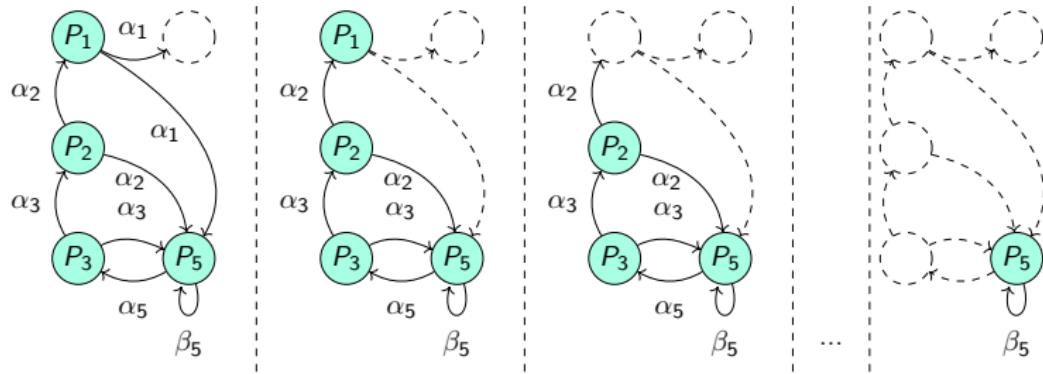
$$\left[\begin{array}{ccc|c} t_{11} & \cdots & t_{1n} & s_1 \\ t_{21} & \cdots & t_{2n} & s_2 \\ \vdots & \ddots & \vdots & \vdots \\ t_{n1} & \cdots & t_{nn} & s_n \end{array} \right]$$

$$\left[\begin{array}{ccc|c} t_{11} & \cdots & t_{1n} & \} \text{ Row Group of } s_1 \\ t_{21} & \cdots & t_{2n} & \} \text{ Row Group of } s_2 \\ \hline t_{31} & \cdots & t_{3n} & \} \text{ Row Group of } s_n \\ \vdots & \ddots & \vdots & \vdots \\ \hline \vdots & \ddots & \vdots & \vdots \\ t_{m1} & \cdots & t_{mn} & \} \text{ Row Group of } s_n \end{array} \right]$$

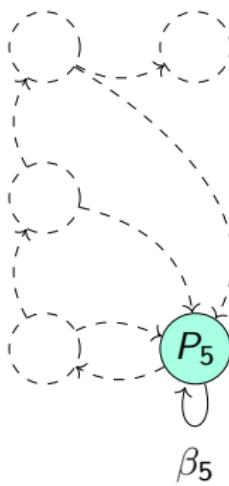
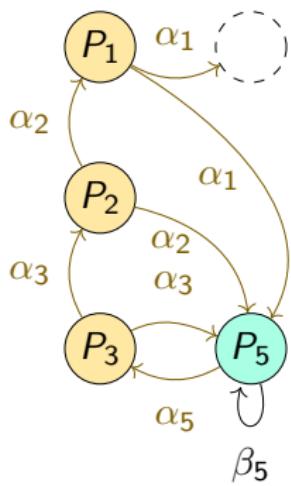
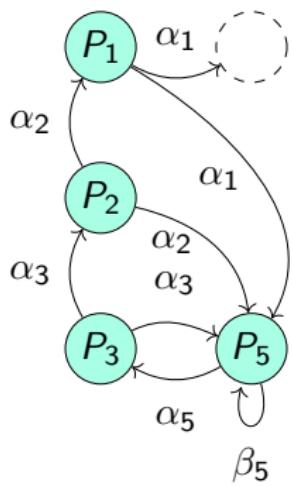
G_{VBA}

G_{EBA}

NAIVE Worst-Case



Random Attractor Removal Optimization



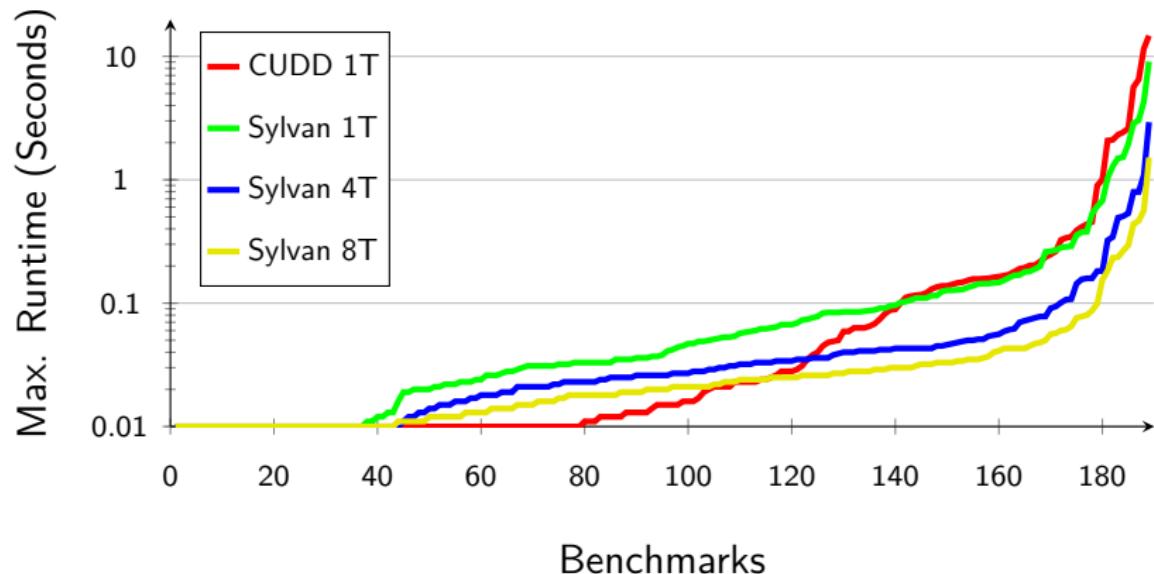
G_{EBA} to G_{VBA} Conversion

Have: $t_{EBA}(\underbrace{x_1, \dots, x_t}_{\text{Source Vertex}}, \underbrace{y_1, \dots, y_u}_{\text{Action}}, \underbrace{x'_1, \dots, x'_t}_{\text{Target Vertex}})$

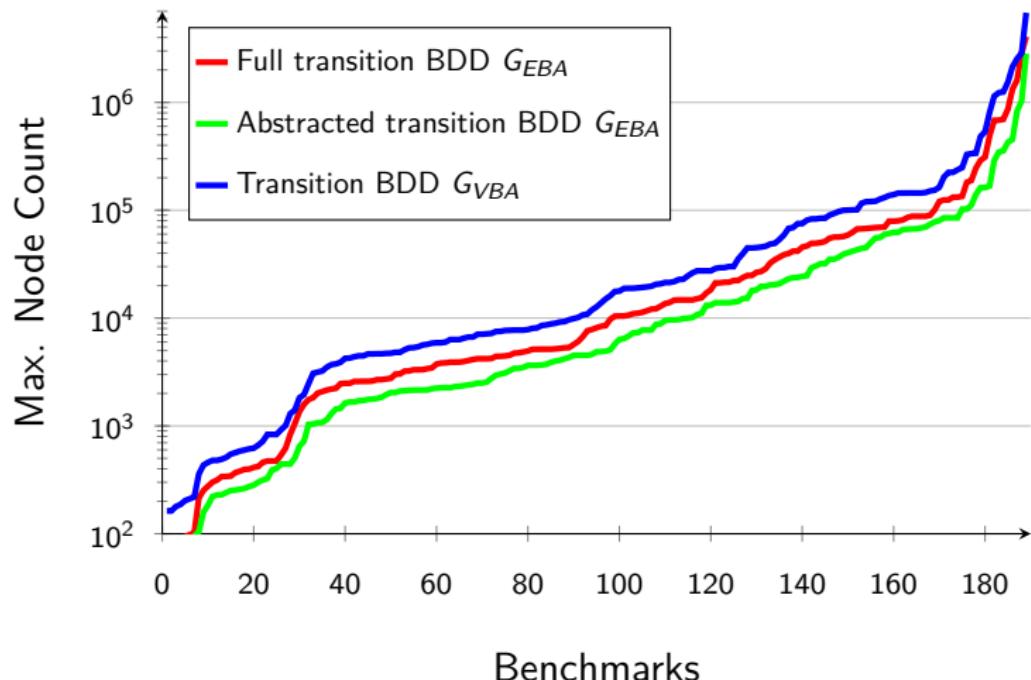
Want: $t_{VBA}(\underbrace{x_1, \dots, \dots, \dots, x_s}_{\text{Source Vertex}}, \underbrace{x'_1, \dots, \dots, \dots, x'_s}_{\text{Target Vertex}})$

Conversion: $t_{VBA}(\underbrace{x_1, \dots, x_t, y_1, \dots, y_u, z}_{\text{Source Vertex}}, \underbrace{x'_1, \dots, x'_t, y'_1, \dots, y'_u, z'}_{\text{Target Vertex}})$

$G_{EBA} \rightarrow G_{VBA}$ Conversion Runtimes



$G_{EBA} \rightarrow G_{VBA}$ Conversion Full Quantile Plot



Benchmarks

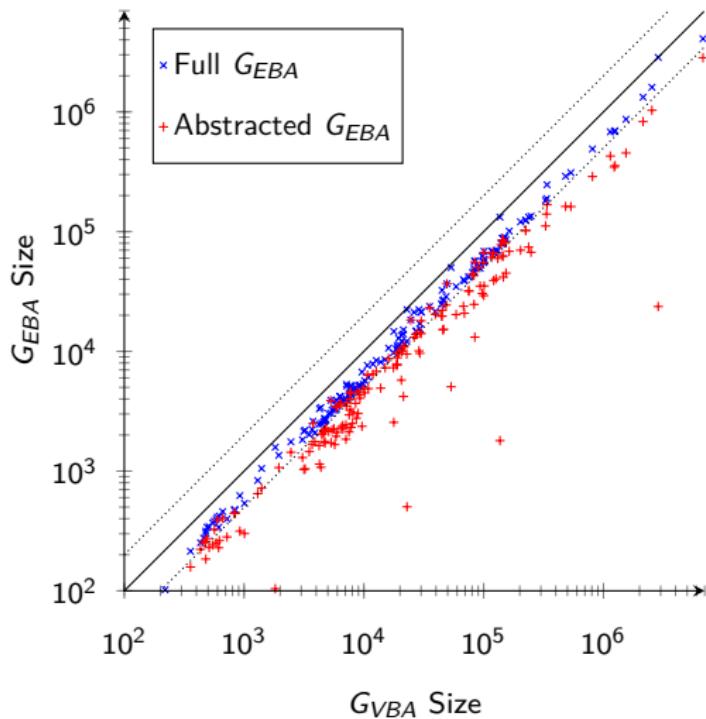
All Transition BDDs

Full: $t_{EBA}(\underbrace{x_1, \dots, x_t}_{\text{Source Vertex}}, \underbrace{y_1, \dots, y_u}_{\text{Action}}, \underbrace{x'_1, \dots, x'_t}_{\text{Target Vertex}})$

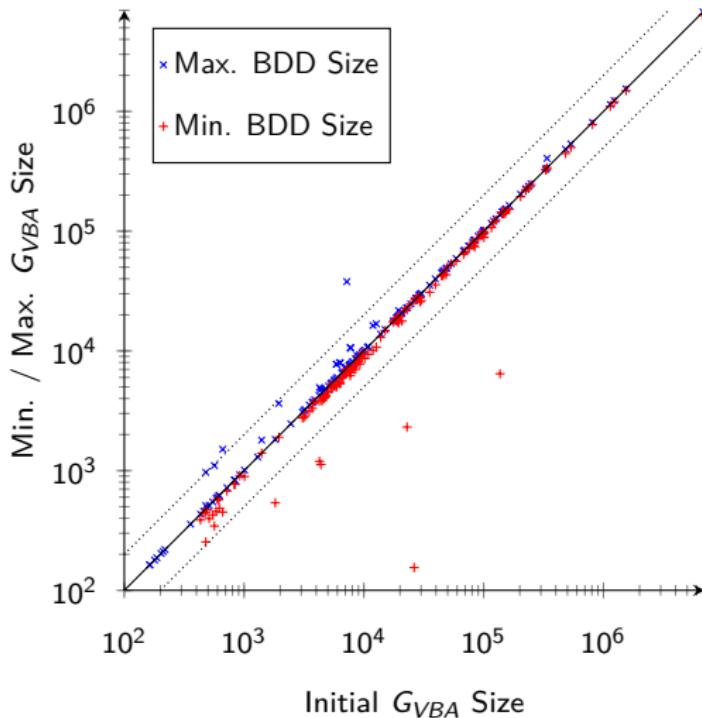
Abstracted: $t_{EBA}(\underbrace{x_1, \dots, x_t}_{\text{Source Vertex}}, \underbrace{x'_1, \dots, x'_t}_{\text{Target Vertex}})$

Converted: $t_{VBA}(\underbrace{x_1, \dots, x_t, y_1, \dots, y_u, z}_{\text{Source Vertex}}, \underbrace{x'_1, \dots, x'_t, y'_1, \dots, y'_u, z'}_{\text{Target Vertex}})$

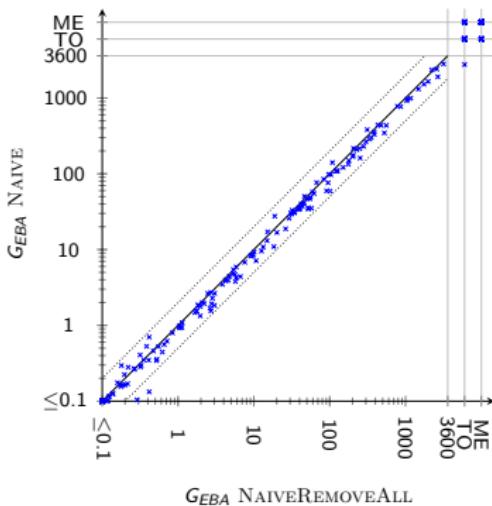
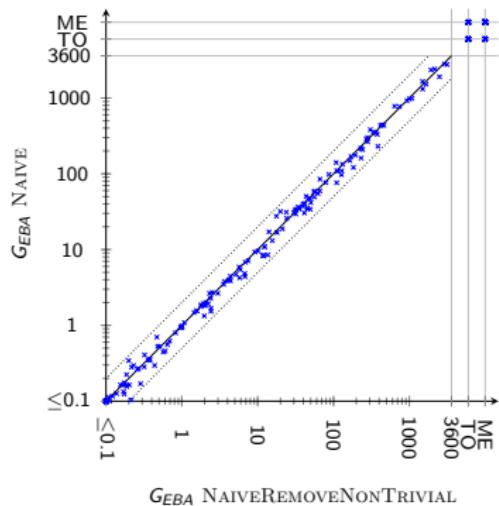
$G_{EBA} \rightarrow G_{VBA}$ Conversion BDD Sizes



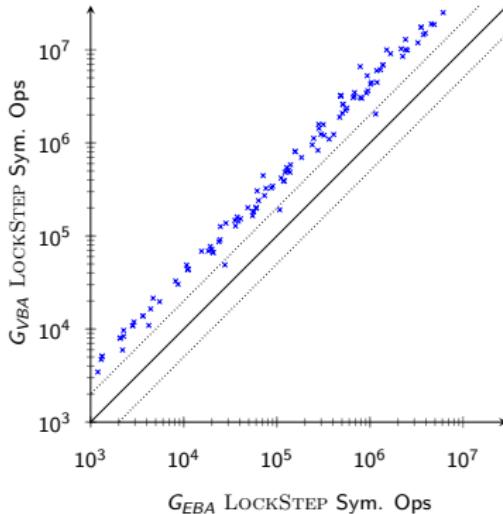
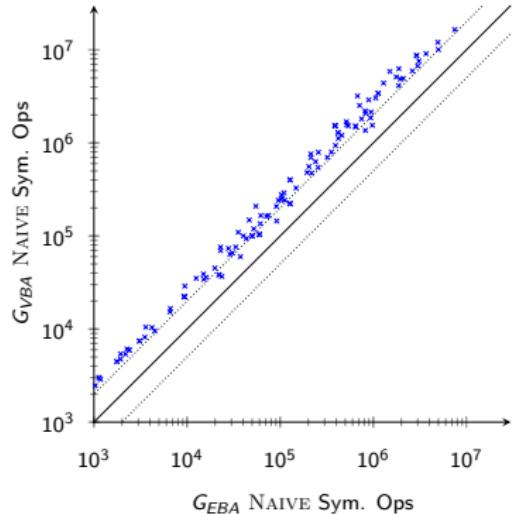
COLLAPSING BDD Sizes



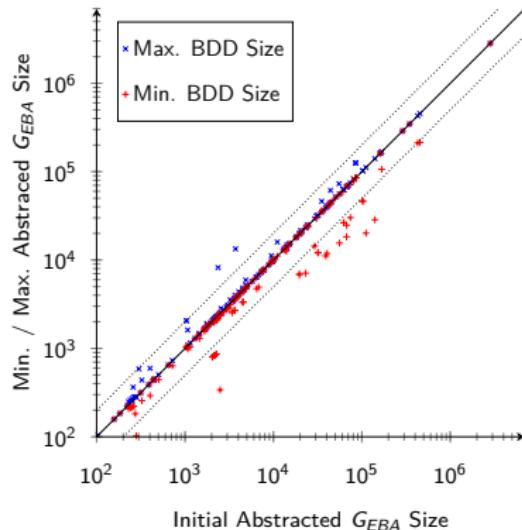
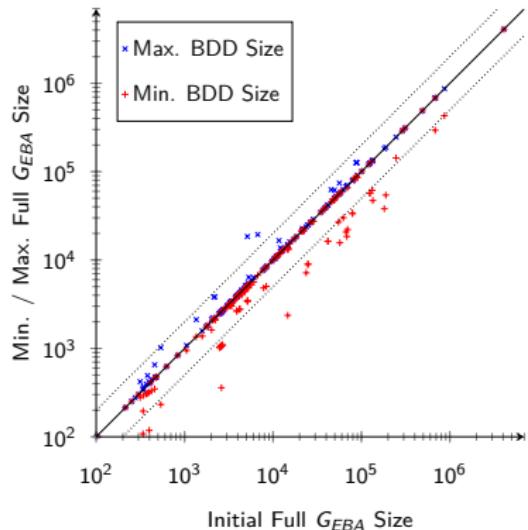
NAIVE Variants Runtime



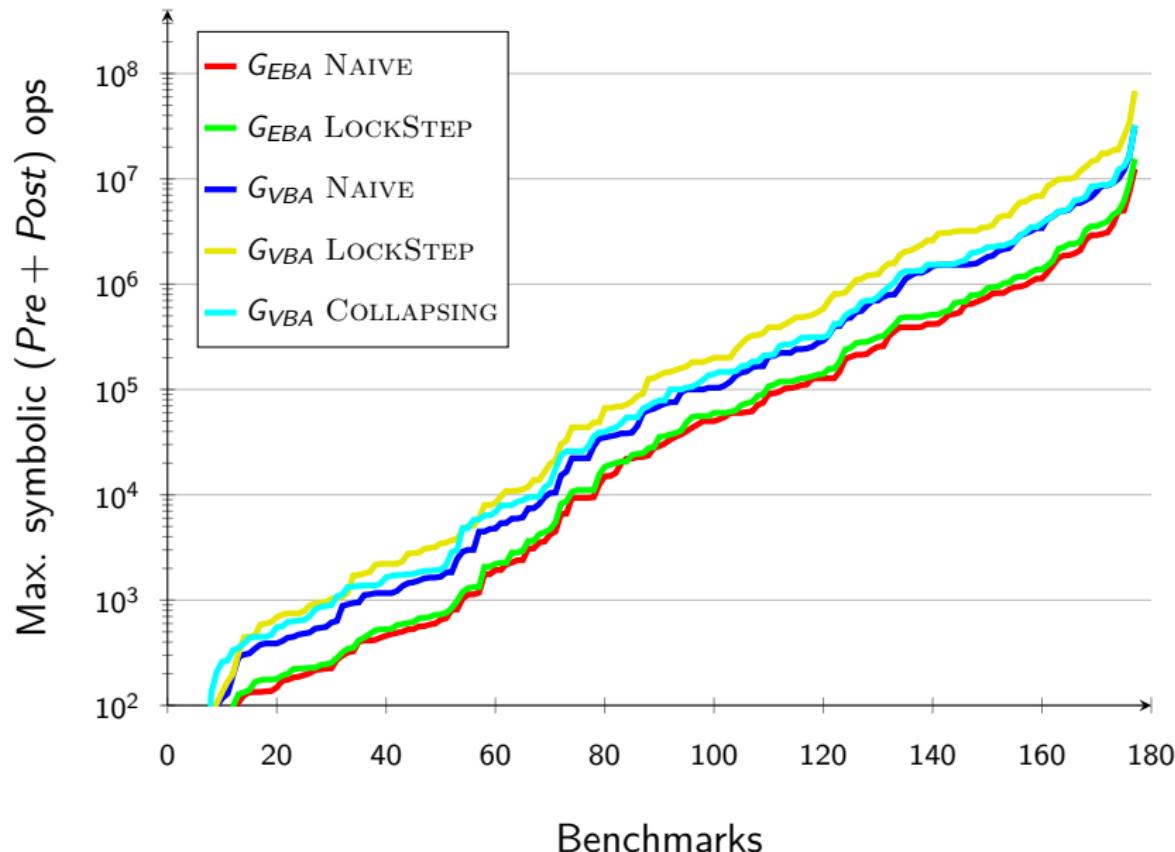
G_{EBA} vs G_{VBA} Sym. Ops



G_{EBA} BDD Sizes



Quantile Plot of Sym. Ops (*Pre + Post*)



Quantile Plot of Sym. Ops (All)

